

Efficient Algorithms for Using Automata in Infinite-State Verification

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Scope of the Thesis

- automata-based FV of parts of real-world programs ...
 - e.g. the Linux OS kernel, C/C++ standard libraries
- ... with complex dynamic data structures,
 - e.g. red-black trees, skip lists
- efficient automata manipulation techniques,
- the use of automata in other applications,
 - e.g. decision procedures of various logics.

- 1 Forest automata-based verification of programs,
- 2 a new decision procedure for a fragment of separation logic,
- 3 ongoing research.

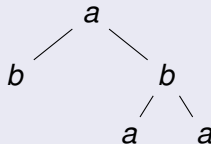
Tree Automata

- An extension of **finite automata** from **words** to **trees**.

a word

aaabbaab

a tree



finite automaton

$q \rightarrow a(p)$

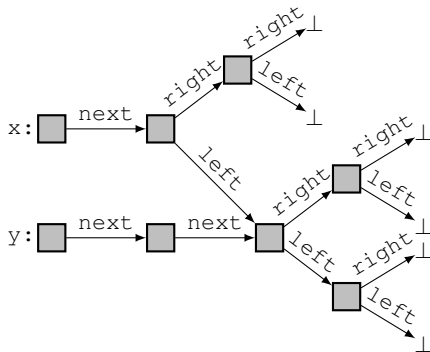
tree automaton

$q \rightarrow a(p_1, \dots, p_n)$

Forest Automata-based Verification

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■ Forest decomposition of a heap

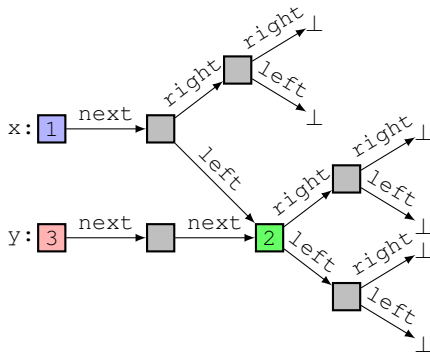


Forest Automata-based Verification

■ Forest decomposition of a heap

- Identify **cut-points**

- nodes referenced:
 - by variables, or
 - multiple times

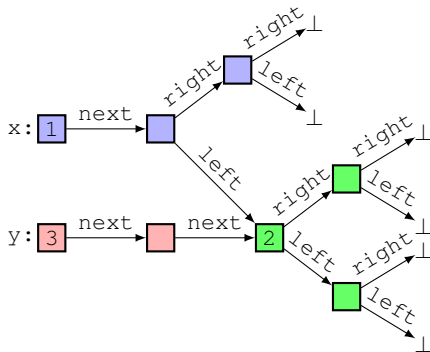


Forest Automata-based Verification

■ Forest decomposition of a heap

- Identify **cut-points**
- Split the heap into **tree components**

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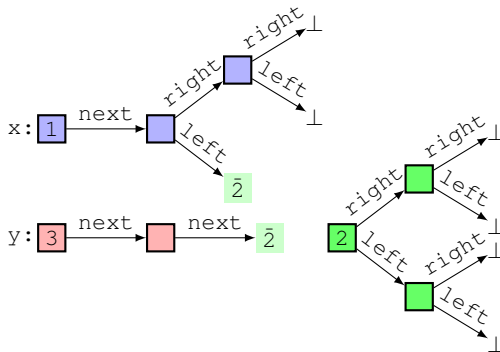


Forest Automata-based Verification

■ Forest decomposition of a heap

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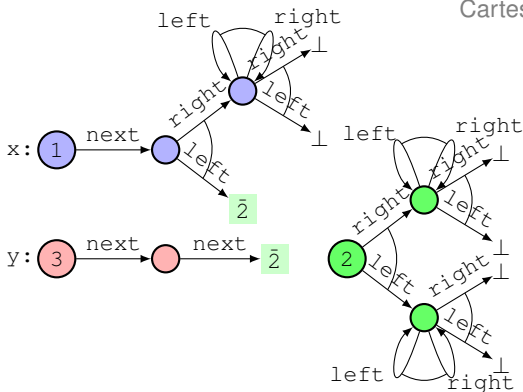


Forest Automata-based Verification

■ Forest decomposition of a heap

- Identify **cut-points**
- Split the heap into **tree components**
 - ▶ **references** are explicit
- Sets of heaps:
 - ▶ **tree automata** to represent sets of tree components
 - ▶ tuple of tree automata $(TA_1, \dots, TA_n) \rightsquigarrow$ **forest automaton** (FA)

- by variables, or
- multiple times



Forest Automata-based Verification

- **Memory safety** of programs manipulating **heap**.
- Uses a **hierarchical** encoding of symbols — called **boxes**,
 - to avoid having infinite tuples of tree components.
- Boxes are discovered automatically,
 - by finding candidate subgraphs to enclose in a box.
- **ARTMC**-based procedure for obtaining a fixpoint.
- The **Forester** tool.

Forest Automata-based Verification (data extension)

- Extending Forester with support for **data-dependent** programs,
 - binary search trees, sorting algorithms, skip lists, ...
- Adding **constraints** \prec_{rr}, \prec_{ra} (**root-root**, **root-all**)
 - locally (inside transitions):

$$q^0 \rightarrow [\text{left}^1(q), \text{right}^2(q)] : 0 \succ_{ra} 1 \wedge 0 \prec_{ra} 2$$

- globally (between a state and a TA):

$$TA_0 \prec_{ar} q \prec_{ra} TA_1$$

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- The need to adjust
 - abstract transformers,
 - inclusion checking,
 - box (un-)folding.
- P. A. Abdulla, L. Holík, B. Jonsson, O. Lengál, C.-Q. Trinh, and T. Vojnar.
[Verification of Heap Manipulating Programs with Ordered Data by Extended Forest Automata.](#)
In *Proc. of ATVA'13*, LNCS 8172.

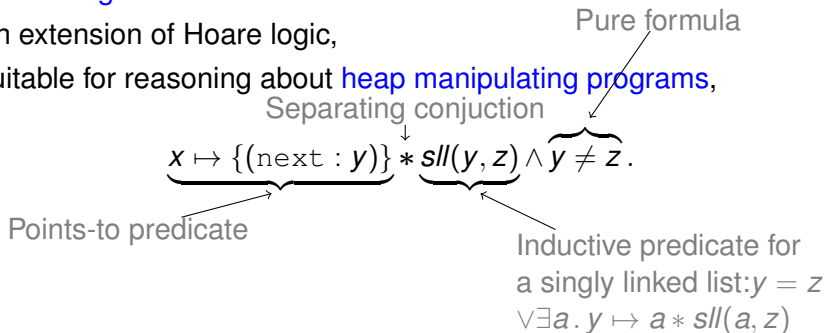
(regular paper)

Deciding a Fragment of Separation Logic

Decision Procedure for a Fragment of Separ. Logic

Separation Logic

- An extension of Hoare logic,
- suitable for reasoning about **heap manipulating programs**,



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Separating conjunction

Pure formula

$$x \mapsto \{(next : y)\} * sll(y, z) \wedge y \neq z.$$

Points-to predicate

Inductive predicate for
a singly linked list: $y = z$

- Checking **entailment**:

$$\varphi \stackrel{?}{\models} \psi$$

$$\forall \exists a. y \mapsto a * sll(a, z)$$

e.g.

$$x \mapsto \{(next : y)\} * sll(y, z) * z \mapsto \{(next : w)\} \stackrel{?}{\models} sll(x, w).$$

- Used e.g. in program analysis for **fixpoint** detection.
- **Undecidable** in general, but there exist **decidable** fragments.

Decision Procedure for a Fragment of Separ. Logic

- The considered fragment is limited to **lists**
 - that are singly/doubly linked,
 - that may be cyclic and/or nested, and
 - that may contain additional head/tail/skip pointers.

Decision Procedure for a Fragment of Separ. Logic

- The decision procedure works as follows:

$$\varphi = \Pi \wedge \varphi_1 * \dots * \varphi_m \stackrel{?}{=} \psi_1 * \dots * \psi_n \wedge \Theta = \psi$$

- 1 **normalize** φ and ψ (adds (dis)equalities, removes emp. segments),
- 2 check that $\Pi \Rightarrow \Theta$,
- 3 for each **points-to** ($x \mapsto \{\dots\}$) in ψ , find a **points-to** in φ ,
- 4 partition the rest of φ , map each class to **inductive predicate** in ψ ,
- 5 for each inductive predicate P in ψ check that $\varphi|_P \models P$ by
 - ① encoding $\varphi|_P$ as a tree $T(\varphi|_P)$,
 - ② encoding P as a tree automaton $TA(P)$, and
 - ③ checking that $T(\varphi|_P) \in L(TA(P))$.

- C. Enea, O. Lengál, M. Sighireanu, T. Vojnar.
Compositional Entailment Checking for a Fragment of Sep. Logic.
Submitted to SAS'14.

(regular paper)

Ongoing Research

Ongoing Research

- 1 A decision procedure for **WSkS**,
 - found a bug in the proposed algorithm, attempting to fix it.
- 2 Algorithms for handling **symbolically represented** automata,
 - computation of **simulation**, **language inclusion**.
- 3 Verification of **concurrent** programs with complex data structures
 - based on forest automata.
- 4 **Adjustable abstraction** in Forester,
 - counterexample-guided abstraction refinement.
- 5 Efficient **language inclusion testing** for tree automata,
 - based on **bisimulation up-to congruence**.

Recent Publications

■ Published:

- L. Holík, O. Lengál, A. Rogalewicz, J. Šimáček, and T. Vojnar.
Fully Automated Shape Analysis Based on Forest Automata. In
Proc. of CAV'13, LNCS 8044. (regular paper)
- P. A. Abdulla, L. Holík, B. Jonsson, O. L., C.-Q. Trinh, and T. Vojnar.
Verification of Heap Manipulating Programs with Ordered Data by
Extended Forest Automata.
In *Proc. of ATVA'13*, LNCS 8172. (regular paper)

■ Submitted:

- P. A. Abdulla, L. Holík, B. Jonsson, O. L., C.-Q. Trinh, and T. Vojnar.
Verification of Heap Manipulating Programs with Ordered Data by
Extended Forest Automata.
Submitted to Acta Informatica (IF'12 = 0.474). (regular paper)
- C. Enea, O. Lengál, M. Sighireanu, T. Vojnar.
Compositional Entailment Checking for a Fragment of Sep. Logic.
Submitted to SAS'14. (regular paper)

Previous Publications

■ Published:

- L. Holík, O. Lengál, J. Šimáček, and T. Vojnar.
Efficient Inclusion Checking on Explicit and Semi-Symbolic TA.
In *Proc. of ATVA'11*, LNCS 6996. (regular paper)
- O. Lengál, J. Šimáček, and T. Vojnar.
VATA: A Library for Efficient Manipulation of Non-Deterministic TA.
In *Proc. of TACAS'12*, LNCS 7214. (regular tool paper: 16 pp)
- O. Lengál.
An Efficient Finite Tree Automata Library.
Lambert Academic Publishing. 2012. (monograph)

■ Products:

- O. Lengál, J. Šimáček, and T. Vojnar.
VATA: A Library for Efficient Manipulation of Non-Deterministic TA.
FIT BUT, 2012. (software)